

# Relations

one CS application: relational databases

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Questions about data stored in relational databases can be posed precisely using the language of relations.

SQL (structured query language) is one language implementing this

Def A binary relation  $R$  on sets  $A, B$  is a subset  $R \subseteq A \times B$ .

Recall:  $A \times B = \{ \langle a, b \rangle : a \in A, b \in B \}$

we write  $(x, y) \in R$  as  $x R y \leftarrow$   
 $(x, y) \notin R$  as  $x \not R y$

examples:  $R_1$

① "is (blood) related to" is a binary relation on people.

Let  $P$  be the set of all people.

"is related to" is  $\{ \langle x, y \rangle : x \in P, y \in P, x \text{ related to } y \}$

$\langle \text{Serena Williams, Venus Williams} \rangle \in R_1$

$\langle \text{Lucy Williams, Venus Williams} \rangle \notin R_1$

②  $<$  on  $A = \{1, 2, 3, 4\}$

$< = \{ \langle 1, 2 \rangle, \langle 1, 3 \rangle, \langle 1, 4 \rangle, \langle 2, 3 \rangle, \langle 2, 4 \rangle, \langle 3, 4 \rangle \}$

$1 < 2$  but  $3 \not< 2$

③ Let  $f: A \rightarrow B$  be a function

$\{ \langle a, f(a) \rangle : a \in A \} \subseteq A \times B$ , so it is a relation

Q Is the converse true?

$\{ \langle x, y \rangle : x \in A, y \in B \} \stackrel{?}{=} \Rightarrow f: A \rightarrow B$  s.t.  $f(x) = y$  is a function?

④ Let  $A = \text{months}$ ,  $B = \text{numbers of days}$

$\{ \langle \text{Jan}, 31 \rangle, \langle \text{Feb}, 28 \rangle, \langle \text{Feb}, 29 \rangle, \langle \text{Mar}, 31 \rangle, \dots \}$

a relation indicating the # days in a month.

Jan 31  
Feb 28  
Feb 29  
Mar 31  
Apr 30  
⋮

Jan → 31  
Feb → 28  
Feb → 29  
Mar →

⑤  $A = \{1, 2, 3, 4\}$   $R_5$

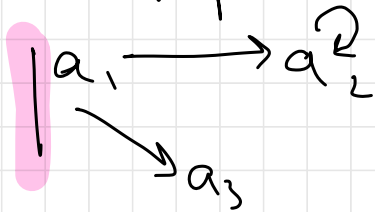
$$R_5 = \{ \langle 1, 1 \rangle, \langle 1, 3 \rangle, \langle 3, 1 \rangle, \langle 3, 3 \rangle, \langle 2, 4 \rangle, \langle 2, 2 \rangle, \langle 4, 2 \rangle, \langle 4, 4 \rangle \}$$

⑥ prereq chart

### Properties of Relations

Let  $R \subseteq A \times A$ , so  $R$  is a relation on  $A$

let's represent  $R$  as a graph:



$$R = \{ \langle a_1, a_2 \rangle, \langle a_2, a_2 \rangle, \langle a_1, a_3 \rangle \}$$

$R$  is reflexive if  $\forall a \in A : a R a$

all nodes have self-loops

$R$  is irreflexive if  $\forall a \in A : a \not R a$

no node has a self-loop

Q are all relations on  $A$  either reflexive or irreflexive?

$$\neg [\forall a \in A : a R a] = \exists a \in A : \neg (a R a) \\ = \exists a \in A : a \not R a$$

$R$  is symmetric if

$$\forall a_1, a_2 \in A : a_1 R a_2 \Rightarrow a_2 R a_1$$



whenever we have a forward edge, we also have a backward edge.

$R$  is anti-symmetric if

$$\forall a_1, a_2 \in A : (a_1 R a_2 \wedge a_2 R a_1) \Rightarrow a_1 = a_2$$



never have backwards edges, but self-loops okay.

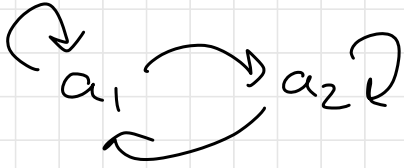
$R$  is transitive if

$$\forall a, b, c \in A : (a R b \wedge b R c) \Rightarrow (a R c)$$

$a \rightarrow b \rightarrow c$  shortcut edges always exist.

Q is  $a_1 \begin{matrix} \rightarrow \\ \leftarrow \end{matrix} a_2$  transitive?

$$(a_1 R a_2 \wedge a_2 R a_1) \Rightarrow (a_1 R a_1)$$



ex relation  $<$  on  $\mathbb{Z}$ .  $1 < 2$

• reflexive:  $\forall a \in \mathbb{Z}: a R a$   $a < a$ ?  
dis

no - proof by counter example.  $1 \not< 1$

• irreflexive:  $\forall a \in \mathbb{Z}: a \not R a$  ( $a \not< a$ )

proof: let  $a \in \mathbb{Z}$ .  $a \not< a$ .  $\square$

• symmetric:  $\forall a, b \in \mathbb{Z}: a < b \Rightarrow b < a$ .

disproof by counter example:  $1 < 2$  but  $2 \not< 1$ .

• antisymmetric: