

Intro to Graphs

$$G = (V, E)$$

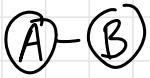
Def An undirected graph $G = (V, E)$ is a non-empty set of vertices / nodes V and a set $E = \{ \{u, v\} : u, v \in V \}$ of edges joining pairs of nodes.

↑
are self-loops allowed?

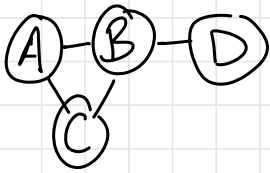
$$\{v\} = \{u, v\} \\ u = v$$

yes!

ex (A) $V = \{A\}$
 $E = \emptyset$



$$V = \{A, B\}$$
$$E = \{ \{A, B\} \} = \{ \{B, A\} \}$$



$$V = \{A, B, C, D\}$$

$$E = \{ \{A, B\}, \{A, C\}, \{B, C\}, \{B, D\} \}$$

non-ex.



(A) all edges need 2 endpoints

real-world examples

- facebook friends

nodes: people

edge: 2 people are friends

- blood relations

Q What property would a relation need to have to be representable as an undirected graph?

ideas:

-irreflexive (no self loops)

-symmetric $a \rightleftarrows b \Leftrightarrow a-b$

(A) (B) is this a valid graph?

$$V = \{A, B\}$$

$$E = \emptyset$$

Def A directed graph $G = (V, E)$ has a set of vertices and edges $E \subseteq V \times V = \{\langle u, v \rangle : u, v \in V\}$ so that edges are directed from one vertex to another.

Q What else looks like $\subseteq V \times V$ relation?

ex (A) $(A) \rightarrow (B) \quad E = \{\langle A, B \rangle\}$

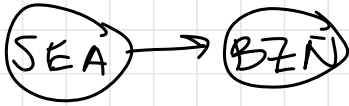
\neq

$(A) \leftarrow (B) \quad E = \{\langle B, A \rangle\}$

- relations are directed graphs
- functions are also directed graphs

real-world exs

- twitter followers
- transportation networks



Def A graph is simple if it contains no parallel edges or self-loops.



note that has no parallel edges.



Example 11.3: Self-loops and parallel edges.

Suppose that we construct a graph to model each of the following phenomena. In which settings do self-loops or parallel edges make sense?

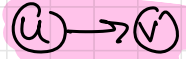
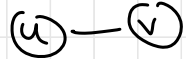
- 1 A social network: nodes correspond to people; (undirected) edges represent friendships.
- 2 The web: nodes correspond to web pages; (directed) edges represent links.
- 3 The flight network for a commercial airline: nodes correspond to airports; (directed) edges denote flights scheduled by the airline in the next month.
- 4 The email network at a college: nodes correspond to students; there is a (directed) edge $\langle u, v \rangle$ if u has sent at least one email to v within the last year.

	<u>Self-loops</u>	<u>parallel edges</u>
social network	probably no	probably no
the web	yes	yes

Flight network probably no yes

email network yes no

Def let $e = \{u, v\}$ or $\langle u, v \rangle$

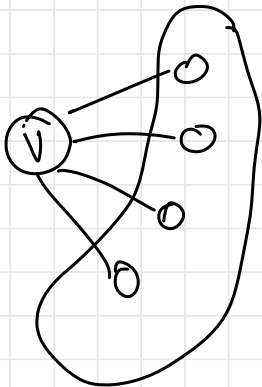


- nodes u, v are adjacent or neighbors
- in a directed graph, v is an out-neighbor of u and u is an in-neighbor of v
- u and v are endpoints of e
- u and v are incident to e ?

(we also say that e is incident to v, u)

let v be a node in an undirected graph.

$$\begin{aligned} \text{degree}(v) &= \text{deg}(v) = d(v) = \# \text{ of neighbors of } v \\ &= \left| \{u \in V : \{v, u\} \in E\} \right| \end{aligned}$$

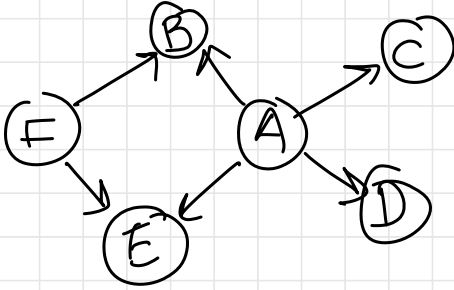


$$\text{deg}(v) = 4$$

for a directed graph, $\text{indeg}(v) = \#$ of in-neighbors of v

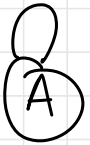
$\text{outdeg}(v) = \# \text{ of}$
 \downarrow
out-neighbors of v

ex



A, B adjacent
D, C not adjacent
A, B are endpoints
of (A, B)

F is an in-neighbor
of B



? \leftarrow let's come back to this