Examples of propositions:

| for ints $n$ | $n(n+1)^{2}$ is even |
| :--- | :--- |
| for ints $n$ | if $n^{2}$ even, men uneven |
| for $x, y \in \mathbb{R}$ | if $x \in \mathbb{Q}, y^{\prime} \in \mathbb{Q}$ men |
| $\sqrt{2}$ is not rational $x y \in \mathbb{Q}$, |  |

In proofs, we've done
$n$ even $\Rightarrow n=2 c$ for $c \in \mathbb{Z} \Rightarrow \cdots$
("implies that")

$$
n x, n y \in \mathbb{Z} \Rightarrow n x n y \in \mathbb{Z}
$$

$\sqrt{2}$ rational $\Rightarrow \cdots \Rightarrow$ false (contradiction)
we can construct compound prop. out of smaller (atomic) prop.
$p$ Ccan't be broken down
if $n$ is integer them $\frac{n(n+1)^{2} \text { even }}{q}$
Syntax vs. Semantics

$\rightarrow$ meaning of a grammatically correct
semfeuce or sentence or
statement
let $p, q$ be prop.

formal semantics (truth table)

if / then
true iff $p$ "forces" $q$ (false if $p$ doesn't it's a promise mat unevever $P_{F} T, q$ also $T$ so $b=7{ }^{\circ}$ is $P_{F}$, when that promise is That is, when $p$ is $T$ and $q$ is $F, p=7 q$ is ex If it rains treen the grass is wet.
 if $p$ then $q$ can also be untten as:
a unenever $P$
$q$ is necessang for $p$
$p$ only if $q /$ scut condition for $q$
lune never $p$ also $q$ lumenever p also q $q$

Deft 2 propositions ave logically equivalent iff their putt tables are the same
$\left[\begin{array}{ccc}p & 1 p & 17 p \\ \hline T & F & T\end{array}\right.$

$$
p \equiv 77 p
$$

pet prop $p$ is satisfiable inf its truth table has af least one T. That is, it's the under at least one fut assignment.
Det A prop. is a tautology iff every row of the tutu fable is $T$
ex $\quad(p=7 q) \vee p$ scraten work

| $p$ | $q$ |
| :---: | :---: |
| $T$ | $T$ |
| $T$ | $F$ |
| $F$ | $F$ |

$$
\begin{array}{cc}
p=q q & (p \Rightarrow q) \vee p \\
\frac{T}{T} & \frac{T}{T} \\
\frac{T}{T} \\
\frac{T}{T}
\end{array}
$$

Q Suppose we have propositions $p, i$, $r$. hallow many vows does the tu th table 8. $\left\{\frac{n}{1}, F\right\}$ general, $2^{n}$. One for seals of

$$
\{T, F\}^{n}=\underbrace{\{\tau, F\} \times\{\tau, F\} \times \cdots \times\{T, F\}}_{n \text { times }}
$$

for $n=3$,

$$
\{T, F\} \times\{T, F\} \times\{T, F\}=\{\langle T, T, T\rangle,\langle T, T, F\rangle, \ldots\}
$$

A drill question: how many rows does the truth table for $a \Rightarrow(b v(c \wedge \neg a))$ have?

$$
2^{3}=8
$$

truth table for

| $a$ | $b$ | $c$ | $\neg a$ | $c \wedge \neg a$ | $b v(c \wedge \neg a)$ | $a \leadsto)(b v(c \wedge \neg))$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $T$ | $T$ | $T$ |  |  |  |  |
| $T$ | $T$ | $F$ |  |  |  |  |
| $T$ | $E$ | $T$ |  |  |  |  |
| $T$ | $F$ | $F$ |  |  |  |  |
| $F$ | $T$ | $T$ |  |  |  |  |
| $F$ | $T$ | $F$ |  |  |  |  |
| $F$ | $F$ | $T$ |  |  |  |  |
| $F$ | $F$ | $F$ |  |  |  |  |
|  |  |  |  |  |  |  |

Precedence Rules
parentheses
1.7

$$
c \wedge 7 a
$$

2. $V, \wedge, \oplus$
3. $\Rightarrow$
4. $\Leftrightarrow$
break firs $L$ to $R$
Def (again) 2 propositions $p, q$ are logicallo equivalent if treir'thenturables ave tue same. we unte $p \equiv q$ if this is true.

| $p$ | $q$ | $\neg p$ | $\neg q$ | $p \vee q$ | $\neg(p \vee q)$ | $\neg p \wedge \neg q$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ | $F$ | $F$ | $T$ | $F$ | $F$ |
| $T$ | $F$ | $F$ | $T$ | $T$ | $F$ | $F$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $F$ | $F$ |
| $T$ | $T$ | $T$ | $F$ | $T$ | $T$ |  |

$$
\neg(p \vee q) \equiv \neg p \wedge \neg q
$$

De Morgan's Law: $\overline{(A \cup B)}=(\bar{A} \cap \bar{B})$

