Intro to Graphs
Deft An undirected graph $G=(V, E)$ is a non-empty set $v$ of nodes/vertices and a set $E=\{\{u, v\}: u, v \in v\}$ of edges joining pairs of nodes.
ex.
(A)

$$
\begin{aligned}
& V=\{A\} \\
& E=\varnothing
\end{aligned}
$$

(A) $-(B)$

$$
\begin{aligned}
& V=\{A, B\} \\
& E=\{\{A, B\}\}=\{\{B, A\}\}
\end{aligned}
$$



$$
\begin{aligned}
& V=\{A, B, C, D\} \\
& E=\{\{A, B\},\{B, C\},\{C, B\},\{B, D\}\}
\end{aligned}
$$

non-ex.
A- all edges need 2 endpoints
real-wond examples:

- facebook friend "nodes = people rages unenthey ave
- blood-related

Q what property would a relation need to be representable as an undirected graph?
Def $A$ directed graph $G=(V, E)$ has set of vertices and edges $t \subseteq V \times V=\{\langle u, v\rangle: u, v \in v\}$ so that edges are directed from one vertex
to another.

$$
\begin{array}{r}
(A) \rightarrow(B) \quad E=\{\langle A, B\rangle\} \\
\neq f \\
(A) \leftarrow(B) \quad E=\{\langle B, A\rangle\}
\end{array}
$$

- relations are directed graphs
- functions ave directed graphs real-wond examples:
- twitter followers
- transportation networks
$S E A \rightarrow B Z N$
Deft $A$ graph is simple if it contains no parallel edges or self-loops.
Parallel edges:

$\rightarrow$ (B) or
self-loops: A) or AR
Example 11.3: Self-loops and parallel edges.
Suppose that we construct a graph to model each of the following phenomena. In which settings do self-loops or parallel edges make sense?

1 A social network: nodes correspond to people; (undirected) edges represent friendships.
2 The web: nodes correspond to web pages; (directed) edges represent links.
3 The flight network for a commercial airline: nodes correspond to airports; (directed) edges denote flights scheduled by the airline in the next month.
4 The email network at a college: nodes correspond to students; there is a (directed) edge $\langle u, v\rangle$ if $u$ has sent at least one email to $v$ within the last year.

Deft
ut $e=\{u, u\}$ or $e=\langle u, v\rangle$

- nodes u,v are adjacent or neighbors
- in a directed graph, $v$ is an out-neignbor of $u$ and $u$ is an in-neignibor of $v$
- $u$ and $v$ are tree endpoints of $e$
- $u$ and $v$ are incident to $e$
let $v$ be a node.

$$
\begin{aligned}
\operatorname{degree}(v)=\operatorname{deg}(v)=\operatorname{d}(v) & =\# \text { neighbors of } v \\
& =|\{u \in v:\{v, u\} \in E\}|
\end{aligned}
$$


for a directed graph,

$$
\begin{aligned}
& \text { indeg }(v)=\text { \# of in-neignbors } \\
& \text { out } \operatorname{deg}(v)=\# \text { of out-neismbors }
\end{aligned}
$$

ex


A,B adjacent $D, C$ not adjacent $A, B$ ave the endpoints of edge $\langle A, B\rangle$
$A, E$ are incident to edge $\langle A, E\rangle$
$F$ is an in-neignbor of $B$
$E$ is an out-neignbor of $A$

