Examples of pRopositions:
for ints $n$ for ints $n$ for neal's $x, y$
$n(n+1)^{2}$ is even if $n^{2}$ even, then $n$ even if $x \in \mathbb{Q}$ and $y \in \mathbb{Q}$ then $x y \in \mathbb{Q}$
$\sqrt{2}$ is not rational
In proofs, we 've done things like:
$n$ even $\Rightarrow n=2 c$ for $\in \mathbb{Z}$
("implies that")

$$
n x, n y \in \mathbb{Z} \Rightarrow n_{x} n_{y} \in \mathbb{Z}
$$

$\sqrt{2}$ rational $\Rightarrow \cdots \Rightarrow \cdots \Rightarrow$ false (contradiction)
$n$ int $\Rightarrow n$ is even or odd
we can consturct compound propositions out of smaller (atomic) propositions.
¿cau't be broken down any smaller
Sytax vs. Semantics
$\rightarrow$ grammatically $\rightarrow$ meaning of a dramatically correct correct sentence or (for given language) statement
let $p, q$ be propositions.

formal semantics

| $p$ | $q$ | $p \wedge q$ | $p \vee q$ | $\neg p$ | $p \Rightarrow q$ | $p \ll>$ | $p \oplus q$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ | $T$ | $T$ | $F$ | $T$ | $T$ | $F$ |
| $T$ | $F$ | $F$ | $T$ | $F$ | $F$ | $F$ | $T$ |
| $F$ | $T$ | $F$ | $T$ | $T$ | $T$ | $F$ | $T$ |
| $F$ | $F$ | $F$ | $F$ | $T$ | $T$ | $T$ | $F$ |

$2 \times$ is even and 3 is odd $T$
2 is even and 4 is odd $F$
1 is even and 3 is odd $F$
3 is even and 2 is odd $F$
2 is even or 3 is odd $T$
2 is even or
1 is even or 4 is odd $T$
3 is even or 3 is odd $T$
2 is odd $F$
not $(2$ is even) $F$
not $(2$ isodd) $T$
if/tren
true iff $p$ "forces" $q$
it's a prom ise mat whenever $p T$, $q$ also $T$ so $p \Rightarrow q$ is $F$ unen that promise is broken. That is, unen $p$ is $T$ and $q$ is $F$.
ex if it rains then the grass is wet. wren is this a lie? rains grasswet

| $T$ | $T$ | $\checkmark$ |
| :--- | :--- | :--- |
| $T$ | $F$ | $x$ |
| $F$ | $F$ | $\checkmark$ |

if $p$ tron $q$ can also be written as:
$q$ whenever $p$
F is necessany for $p$
$b$ only if a
$p$ is a sufficient condition for $p$
unevever $p$ also $q$
p implies q
Q suppose we have propositions $p, q, r$.
how many rows does the truth table have?
$2^{n}$. one for earn of $\{T, F\}^{n}$. (recall set notation)

$$
\begin{aligned}
& \{T, F\} \times\{T, F\}=\{\langle T, T \geqslant,\langle T, F\rangle,\langle F, T\rangle,\langle F, F\rangle\} \\
& \{T, F\}^{\}}=\{\langle T, T, T\rangle,\langle T, T, F\rangle, \ldots\}
\end{aligned}
$$

Det 2 propositions are logically equivalent, whiten ミ, iff their tuith tables are the same.


| $p$ | $q$ | $p \Rightarrow q$ | $q \Rightarrow p$ | $\neg p \vee q$ | $\neg q \Rightarrow-p$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ | $T$ | $T$ | $T$ | $T$ |
| $T$ | $F$ | $F$ | $T$ | $F$ | $F$ |
| $F$ | $T$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $T$ | $T$ | $T$ | $T$ |

$$
\begin{aligned}
& (p \Rightarrow \neg q) \equiv(\neg p \vee q) \\
& (p \Rightarrow \neg q) \\
& (p \Rightarrow \neg q) \neq(\neg q \Rightarrow \neg p) \\
& \neg(p \vee q) \equiv(\neg p) \wedge(\neg q) \\
& (\overline{A \cup B})=(\bar{A} \cap \bar{B})
\end{aligned}
$$

De Morgan's Law

Precedence rules

1. 7
2. $v_{1} \wedge_{1} \oplus$
3. $=7$
4. $<>$

Deft Prop $p$ is satisfiable iff it is the under at least one tret assignment.
that is, at least one vow of the turn table evaluates to $T$.
ex

$$
p_{p} \stackrel{\rightharpoonup q}{ }
$$

Deft A prop is a tautology if every row of the truman table is the.

$$
\text { ex }(p \Rightarrow q) \wedge p
$$

