

Practice Quiz 2

Names \_\_\_\_\_

Problem 1 (20 points)

In this problem, you will prove that  $2 \log n = O(\log n)$  using the definition of big O. Follow the three steps carefully.

(5 points) Write down the definition of big O.

(10 points) Give a  $c$  and a  $n_0$  that can be used to prove that  $2 \log n = O(\log n)$ .

(5 points) Explain what it would mean for this  $c$  and  $n_0$  to work in a proof that  $2 \log n = O(\log n)$ , and *very briefly* explain why they do (write one sentence, draw a graph, etc).

Problem 2 (20 points)

Complete the proof that  $\sum_{i=1}^n (2i - 1) = n^2$  using mathematical induction by filling in the following. Note that  $\sum_{i=1}^n (2i - 1) = 1 + 3 + 5 \cdots + (2n - 1)$ . Each underline is worth one point. The proof of the base case is worth two points. The proof of the inductive case is worth eight points.

*Proof.* We prove that  $\sum_{i=1}^n (2i - 1) = n^2$ .

First, let  $P(n)$  be the predicate that \_\_\_\_\_.

We prove that \_\_\_\_\_ (something to do with  $P$ ) using mathematical induction over  $n$ .

*Base case.* We show that \_\_\_\_\_ (something to do with  $P$ ).

*Inductive case.* We show that \_\_\_\_\_ (something to do with  $P$ ).

Assume \_\_\_\_\_ (something to do with  $P$ ). That is, \_\_\_\_\_  
(translating the previous blank using the formula, aka, the inductive hypothesis, or IH.).

Because we proved both \_\_\_\_\_ and \_\_\_\_\_, by the principle of mathematical induction,  $\forall n \geq 0 : P(n)$ . □

Problem 3 ( 40 points)

- (a) Give an example of a function that is **not** one-to-one.
- (b) Suppose  $f : A \rightarrow B$  is both onto and one-to-one. How do  $|A|$  and  $|B|$  compare?
- (c) Let  $f : \mathbb{Z} \rightarrow \mathbb{Z}$  be defined as  $f(n) = 2n$ . What is the codomain of  $f$ ? What is the range of  $f$ ?
- (d) Suppose we have a recurrence relation describing the runtime of a recursive algorithm as follows:  
 $T(0) = 5$  and  $T(n) = T(n - 1) + 3$ . What is  $T(2)$ ?
- (e) Is  $n = O(\log n)$ ?
- (f) Is  $\log n = O(n)$ ?
- (g) Fill in the rest of the recursive definition of a linked list.  
A linked list is either
- (1) An empty list  $\langle \rangle$ , or
  - (2)  $\langle x, L \rangle$  where  $x$  is a data element and  $L$  is \_\_\_\_\_

Problem 4 (20 points)

- (a) For the following algorithm give a proposed function representing the number of primitive operations for the algorithm in terms of the input size, addressing each line and/or loop of the algorithm. You do not need to be precise counting constant numbers of primitive operations (e.g., figuring out exactly how many primitive operations a line does). However, you should be precise about how many times a loop runs.

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**Algorithm 1**

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1: for  $i = 1$  to  $n \cdot n$  do
2:   if  $3|i$  then
3:     for  $j = 1$  to  $n$  do
4:        $x = x + 1$ 
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Algorithm 1 takes  $f(n) =$  \_\_\_\_\_ primitive operations.

- (b) For the  $f(n)$  you gave in (a), give the “tightest” (or asymptotically smallest)  $g(n)$  such that  $f(n) = O(g(n))$ .

$f(n) = O($  \_\_\_\_\_  $)$