INSTRUCTOR: LUCIA WILLIAMS CSCI 332: ADVANCED ALGORITHMS & DATA STRUCTURES

After you sit down, please fold your paper hot dog style and write:

‣What you'd like to be called **‣**Your hometown

‣Your pronouns

‣Your major/concentration

‣A fun fact about you

computer Seattle, WA LUCY (or Professor Williams) I have two huge she/her dogs

Introduce yourself to your neighbors!

" An algorithm is a finite, definite, effective procedure, with some input and some output. "

 — Donald Knuth

"Algorithmic problems form the heart of computer science, but they rarely arrive as cleanly packaged, mathematically precise questions. Rather, they tend to come bundled together with lots of messy, application-specific detail, some of it essential, some of it extraneous."

 — Kleinberg & Tardos

What were the focuses of CSCI 232?

CSCI 232. Implementation and consumption of classic algorithms.

独Fundamental data structures (arrays, stacks, queues, etc.).

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- Sorting.
- Searching.
- Graph algorithms.
- String processing.
- Compression.

```
private static void sort(double[] a, int lo, int hi) { 
  if (hi \leq 10) return;
  int It = Io, gt = hi;
 int i = Io;while (i \leq g) {
    if (a[i] < a[lo]) swap(a, [t++, i++);
    else if (a[i] > a[lo]) swap(a, i, gt-);
     else i++; 
   }
  sort(a, lo, lt - 1);
  sort(a, gt + 1, hi);}
```
Emphasizes critical thinking, problem-solving, and code.

CSCI 332. Design and analysis of algorithms.

• Finding computational problems in the real world.

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- Duality.
- Data structures.
- Intractability.

Emphasizes critical thinking, problem-solving, and both open-ended problems and rigorous analysis. ⁶ *" Algorithms are the life-blood of computer science… the common denominator that underlies and unifies the different branches. " — Donald Knuth*

Why study algorithms?

 $\ddot{\cdot}$

Internet. Web search, packet routing, distributed file sharing, ... Biology. Human genome project, protein folding, … Computers. Circuit layout, databases, caching, networking, compilers, … Computer graphics. Movies, video games, virtual reality, … Security. Cell phones, e-commerce, voting machines, … Multimedia. MP3, JPG, DivX, HDTV, face recognition, … Social networks. Recommendations, news feeds, advertisements, … Physics. Particle collision simulation, *n*-body simulation, …

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We emphasize algorithms and techniques that are useful in practice.

In table groups, try to complete the syllabus quiz. Some of the questions are openended and may not have one single answer!

If your group comes up with a question you can't answer (not necessarily one on the quiz), post it in #questions in Discord.

How to match? What should we think about when designing an algorithm for this problem?

Given:

* a set of preferences among hospitals and med-school students

hospitals' preference lists

students' preference lists

* a matching of hospitals to students

{ A–Z, B-Y, C-X }

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least favorite favorite **1st 2nd 3rd** Xavier Boston Atlanta Chicago Yolanda Atlanta Boston Chicago **Zeus** Atlanta Boston Chicago least favorite

students' preference lists

* a matching of hospitals to students

{ A–Z, B-Y, C-X }

With your table group, give at least two *measurable* criterion for a "good" matching.

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\n
$$
\bigvee_{\begin{array}{c}\n\searrow \\
\searrow\n\end{array}}
$$

The score is the sum of the ranks for every pair. Smaller scores are better.

Worksheet

You have 15 minutes. Ask for help if needed.

For *n* hospitals/students, how many unique matchings?

Algorithm to finding matching with best score?

brute force - try all $H_{5} \Box$

Runtime?

 M

-small example

-

Goal. Given a set of preferences among hospitals and med-school students, design a self-reinforcing admissions process. Matching med-
Goal. Given a set
a self-reinforcing a

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Unstable pair. Hospital *h* and student *s* form an unstable pair if both:

- h prefers s to $\mathcal{A}\mathcal{B}$ of its admitted students. pspita
BBC
	- *s* prefers *h* to assigned hospital.

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- *s* prefers *h* to assigned hospital.

Stable assignment. Assignment with no unstable pairs.

• Individual self-interest prevents any hospital–student side deal.

Stable matching problem: input

Input. A set of *n* hospitals *H* and a set of *n* students *S*.

one student per hospital (for now)

 \blacktriangledown

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 \cdot Each hospital h ∈ H ranks students.

one student per hospital (for now)

hospitals' preference lists

Stable matching problem: input

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- Each hospital $h \in H$ ranks students.
- Each student $s \in S$ ranks hospitals.

one student per hospital (for now)

 $\boldsymbol{\mathcal{N}}$

hospitals' preference lists

Stable matching problem: output

Def. A set $M \subseteq H \times S$ is a matching if and only if:

a perfect matching $M = \{ A-Z, B-Y, C-X \}$

Stable matching problem: output

Def. A set $M \subseteq H \times S$ is a matching if and only if:

- Each hospital $h \in H$ appears in at most one pair of M .
- Each student $s \in S$ appears in at most one pair of M .

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- Each hospital $h \in H$ appears in at most one pair of M .
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Def. A matching M is perfect if $|M| = |H| = |S| = n$.

a perfect matching $M = \{ A-Z, B-Y, C-X \}$

Unstable pair

Def. Given a perfect matching *M*, hospital *h* and student *s* form an unstable pair if both:

- *h* prefers *s* to matched student.
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Key point. An unstable pair *h–s* could each improve by joint action.

A-Y is an unstable pair for matching $M = \{A-Z, B-Y, C-X\}$

On your own, think about…

4. None of the above.

On your own, think about…

Which pair is unstable in the matching { A-X, B-Z, C-Y }?

- 1. A–Y.
- 2. B–X.
- 3. B–Z.
- 4. None of the above.

B-X is an unstable pair

Stable matching problem

Def. A stable matching is a perfect matching with no unstable pairs.

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Do you see any potential issues with using Stable Matching to solve the med student to hospital matching problem?

- 2*n* people; each person ranks others from 1 to $2n 1$.
- Assign roommate pairs so that no unstable pairs.

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5 minute break

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Let's vote

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5 minute break

Gale-Shapley deferred acceptance algorithm

An intuitive method that guarantees to find a stable matching.

GALE–SHAPLEY (preference lists for hospitals and students)

INITIALIZE M to empty matching.

WHILE (some hospital h is unmatched and hasn't proposed to every student)

 $s \leftarrow$ first student on h's list to whom h has not yet proposed.

IF (s is unmatched)

Add h -s to matching M.

```
ELSE IF (s prefers h to current partner h')
```
Replace h' –s with h –s in matching M.

```
ELSE
```
s rejects h .

RETURN stable matching M.

Gale–Shapley demo

hospitals' preference lists

initialize M

We enter the while loop. How many valid first steps are there?

Gale–Shapley demo

hospitals' preference lists

Atlanta proposes to ????

Atlanta proposes to Wayne

students' preference lists

Atlanta proposes to Wayne

Wayne accepts (since previously unmatched)

Boston proposes to Yolanda

students' preference lists

Boston proposes to Yolanda

Yolanda accepts (since previously unmatched)

Chicago proposes to Wayne

What happens?

Chicago proposes to Wayne

Gale–Shapley demo

hospitals' preference lists

students' preference lists

Chicago proposes to Wayne

Wayne accepts (and renounces Atlanta)

Gale–Shapley demo

hospitals' preference lists

Atlanta proposes to Val

students' preference lists

Atlanta proposes to Val

Val accepts (since previously unmatched)

Gale–Shapley demo

hospitals' preference lists

Detroit proposes to Val

hospitals' preference lists

students' preference lists

Detroit proposes to Val

Val rejects (since she prefers Atlanta)

hospitals' preference lists

Detroit proposes to Yolanda

hospitals' preference lists

students' preference lists

Detroit proposes to Yolanda

Yolanda accepts (and renounces Boston)

hospitals' preference lists

Boston proposes to Wayne

hospitals' preference lists

students' preference lists

Boston proposes to Wayne

Wayne rejects (since he prefers Chicago)

hospitals' preference lists

Boston proposes to Val

hospitals' preference lists

students' preference lists

Boston proposes to Val

Val rejects (since she prefers Atlanta)

hospitals' preference lists

Boston proposes to Xavier

hospitals' preference lists

students' preference lists

Boston proposes to Xavier

Xavier accepts (since previously unmatched)

hospitals' preference lists

El Paso proposes to Wayne

hospitals' preference lists

students' preference lists

El Paso proposes to Wayne

Wayne rejects (since he prefers Chicago)

hospitals' preference lists

El Paso proposes to Yolanda

hospitals' preference lists

students' preference lists

El Paso proposes to Yolanda

Yolanda accepts (and renounces Detroit)

hospitals' preference lists

Detroit proposes to Xavier

hospitals' preference lists

students' preference lists

Detroit proposes to Xavier

Xavier rejects (since he prefers Boston)

hospitals' preference lists

Detroit proposes to Wayne

hospitals' preference lists

students' preference lists

Detroit proposes to Wayne

Wayne rejects (since he prefers Chicago)

hospitals' preference lists

Detroit proposes to Zeus

hospitals' preference lists

students' preference lists

Detroit proposes to Zeus

Zeus accepts (since previously unmatched)

hospitals' preference lists

Can Gale-Shapley ever result in an infinite loop?

1. Yes

2. No

What is the worst-case runtime of Gale-Shapley on an input of size n?

1. log *n*

2. *n*

3. n^2

4. *n*!

Observation 2. Once a student is matched, the student never becomes unmatched; only "trades up."

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	1 st	2 _{nd}	$\sqrt{3}$ rd	-4 th	5 th		1 st	2 _{nd}	3rd	4 th	-5 th
\mathbf{A}	\vee	W	X	Y	Z	V	$\mathsf B$	$\mathsf C$	D	E	A
\mathbf{B}	W	X	Υ	V	Z	W	$\mathbf C$	D	E	A	B
C	\boldsymbol{X}	Y	V	W	Z	X	D	E.	A	B	$\mathsf C$
D	Y	V	W	\boldsymbol{X}	Z	Y	E.	A	B	$\mathsf C$	D
E.	\vee	W	X	Y	Z	Z	\overline{A}	B	$\mathsf C$	D	E.

n(n-1) + 1 proposals

Does any hospital end up with more than one student?

1. Yes

2. No

Proof of correctness: perfect matching

Claim. Gale–Shapley outputs a matching. Pf.

• Hospital proposes only if unmatched. \Rightarrow matched to \leq 1 student

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- ・Student keeps only best hospital. ⇒ matched to ≤ 1 hospital

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	- ・By previous claim, all *n* hospitals get matched.
	- Thus, all *n* students get matched.

Claim. In Gale–Shapley matching *M**, there are no unstable pairs.

Pf. Consider any pair *h*–*s* that is not in *M**.

Gale–Shapley matching M*

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hospitals propose in decreasing order of preference

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⇒ *s* rejected *h* (either right away or later)

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- Always returns a stable matching (cases)

open questions : - hospital optimal ? -

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Why did we focus on stable matching instead of minimum score matching?

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