

what did you notice about polynomial functions vs. exponential functions?

when you double the input  $n$  to a polynomial function  $f = an^b$ , the output increases by a constant.

which constant?  $2^b$

when you double the input  $n$  to an exponential function  $a \cdot b^n$ , the output increases by a function of  $n$ .

which function of  $n$ ?  $b^n$

	$f(n) = n^2$	$f(n) = 2^n$	$f(n) = n^{100}$
$n = 10$	100	1028	
$n = 20$	400	$1028 \cdot 1028$ $\approx 1m!$	
$n = 100$	10,000 $10^5$	$10^{30}$	
$n = 200$	40,000 $10^5$	$10^{60}$	

at what point does  $n^{100}$  start being better than  $2^n$ ?

we say an algorithm is efficient if it is polynomial time.

Is there an efficient alg for stable matching?

Is there an efficient alg for sorting?

Is there a problem for which there is no efficient algorithm?



Why did I have you do the doubling input question?

Can you make any general statements?

$Cn^a$  goes up by  $2^a$  when  $n$  doubles

$Kb^n$  goes up by  $b^n$  when  $n$  doubles

$$\frac{C(2n)^a}{Cn^a} = \frac{2^a \cancel{Cn^a}}{\cancel{Cn^a}} = 2^a$$

$$\frac{Kb^{2n}}{Kb^n} = b^{2n-n} = b^n$$

so one thing we can mean by "scales" is doubling input doesn't ↑ as  $n$  ↑

What about logs?

$\log_b n$

$$\frac{\log_b 2n}{\log_b n} = \frac{\log_b 2}{\log_b n} + \frac{\log_b n}{\log_b n}$$

$$= \frac{\log_b 2}{\log_b n} + \frac{\log_b n}{\log_b n}$$

$$= 1 + \frac{\log_b 2}{\log_b n}$$

as  $n \rightarrow \infty$ , ?



## Asymptotic Bounds

Let  $f(n)$  and  $g(n)$  be functions.

$f(n)$  is  $O(g(n))$  if  $f(n)$  is asymptotically upper bounded by  $g(n)$ .

$$f(n) = O(g(n))$$

more precisely, if  $f(n)$  is upper bounded by a constant multiple of  $g(n)$  for sufficiently large  $n$ .

ex  $f(n) = 10n^2 + 3n + 1000$

$$g(n) = n^2$$

is  $f(n) = O(g(n))$ ?



let's apply def.

① no  $\leftarrow$  first

②

③ no  $\leftarrow$  second

ex  $f_2(n) = n \log n$   
 $g(n) = n^4$

is  $f(n) = O(g(n))$

$f(n)$  is  $\Omega(g(n))$  if  $f(n)$  is asymptotically  
lower bounded by  $g(n)$ .  
more precisely, fill in w/ groups

w/ groups, is  $f_1(n) = \Omega(g(n))$ ?  $f(n) = \Omega(h(n))$ ?

$f(n)$  is  $\Theta(n)$  if  $f(n)$  is  $O(n)$  and  $f(n)$  is  $\Omega(n)$   
don't forget mnemonics

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