

Integer Vertex Cover Linear Program

Input

Graph $G = (V, E)$

Where $V = \{v_1, v_2, \dots, v_n\}$

and $E = \{\{v_i, v_j\} \text{ where } v_i, v_j \in V\}$

Output

smallest set of vertices covering all edges

Linear Program

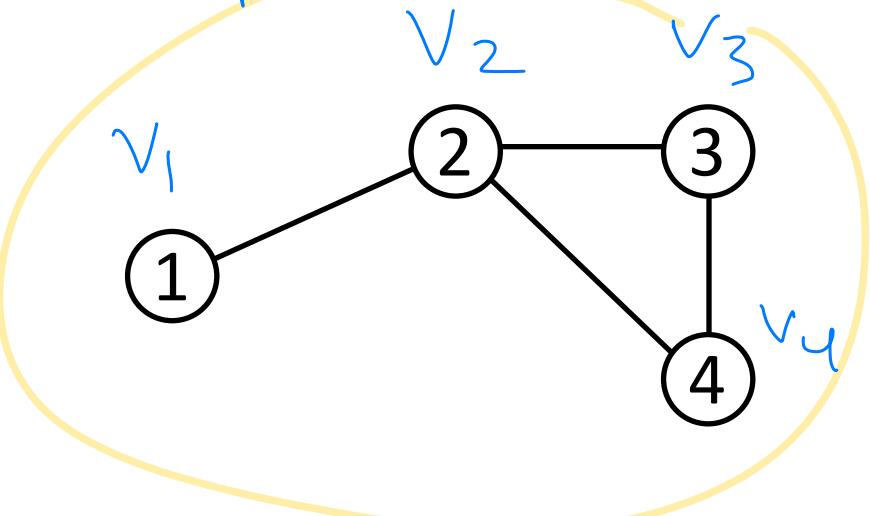
vars: x_i for each vert v_i
 $x_i \in \{0, 1\}$

objective:

$$\min \sum x_i$$

constraints:

for each edge,
sum of endpoints ≥ 1



Vertex Cover ILP

Vertex Cover: Given graph, find the smallest subset of vertices such that every edge in the graph has at least one vertex in the subset.

Objective: $\min \sum_i x_i$

Subject to: $x_i + x_j \geq 1$, for each edge $\{v_i, v_j\}$

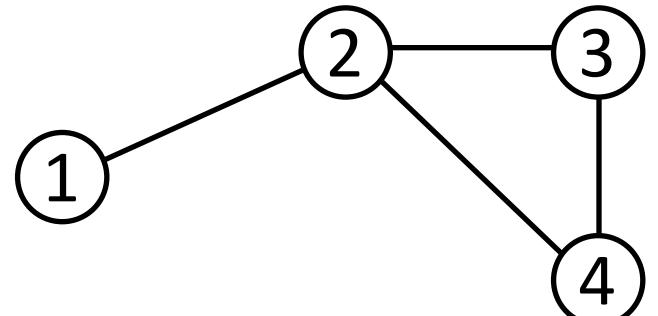
$x_i \in \{0,1\}$, for each vertex i

Example:

Objective: $\min x_1 + x_2 + x_3 + x_4$

Subject to: $x_1 + x_2 \geq 1$

$x_1, x_2, x_3, x_4 \in \{0,1\}$



Vertex Cover ILP

binteger

Vertex Cover: Given graph, find the smallest subset of vertices such that every edge in the graph has at least one vertex in the subset.

Objective: $\min \sum_i x_i$

Subject to: $x_i + x_j \geq 1$, for each edge $\{v_i, v_j\}$

$x_i \in \{0,1\}$, for each vertex i

Example:

Objective: $\min x_1 + x_2 + x_3 + x_4$

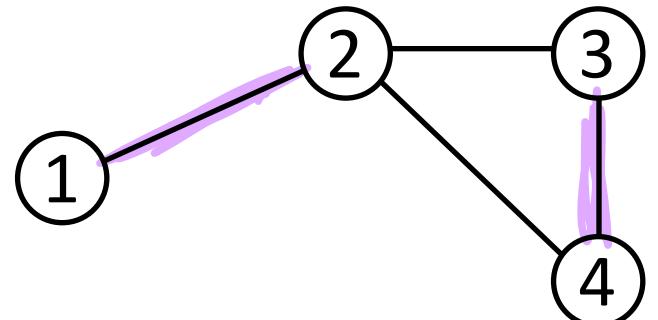
Subject to: $x_1 + x_2 \geq 1$

$x_2 + x_3 \geq 1$

$x_2 + x_4 \geq 1$

$x_3 + x_4 \geq 1$

$x_1, x_2, x_3, x_4 \in \{0,1\}$



Set Cover ILP

Set Cover: Given a universe of elements U and sets S , find the smallest subset of S such that every element in U is in some selected subset.

$$U = \{1, 4, 7, 8, 10\}$$
$$S = \left\{ \begin{array}{l} \{1, 7, 8\}, \{1, 4, 7\}, \\ \{7, 8\}, \{4, 8, 10\} \end{array} \right\}$$

Set Cover ILP

Set Cover: Given a universe of elements U and sets S , find the smallest subset of S such that every element in U is in some selected subset.

Objective: $\min \sum_s x_s$

Subject to: $\sum_{s: u \in s} x_s \geq 1$, for each $u \in U$
 $x_s \in \{0,1\}$, for each set s

$$x_1 + x_2 \geq 1$$

$$U = \{1, 4, 7, 8, 10\}$$
$$S = \left\{ \begin{array}{l} \{1, 7, 8\} \\ \{7, 8\}, \{4, 8, 10\} \end{array} \right\}$$
$$\begin{array}{ll} x_1 = 1 & x_3 = 0 \\ x_2 = 0 & x_4 = 1 \\ x_3 = 1 & x_4 = 0 \end{array}$$

Set Cover ILP

Set Cover: Given a universe of elements U and sets S , find the smallest subset of S such that every element in U is in some selected subset.

Objective:

$$\min \sum_s x_s$$

Subject to:

$$\sum_{s: u \in s} x_s \geq 1, \text{ for each } u \in U$$
$$x_s \in \{0,1\}, \text{ for each set } s$$

Example:

$$\text{Objective: } \min x_1 + x_2 + x_3 + x_4$$

$$\text{Subject to: } x_1 + x_2 \geq 1$$

$$x_2 + x_4 \geq 1$$

$$x_1 + x_2 + x_3 \geq 1$$

$$x_1 + x_3 + x_4 \geq 1$$

$$x_4 \geq 1$$

$$x_1, x_2, x_3, x_4 \in \{0,1\}$$

$$U = \{1, 4, 7, 8, 10\}$$

$$S = \left\{ \begin{array}{l} \{1, 7, 8\}, \{1, 4, 7\}, \\ \{7, 8\}, \{4, 8, 10\} \end{array} \right\}$$

We now have a *poly-time reduction* from Vertex Cover to Set Cover

Integer linear programming
EP

objective, constraints
linear funcs of vars

Vertex Cover and Set Cover are NP-hard

NP-hard = if we can solve in polynomial time, then $P = NP$

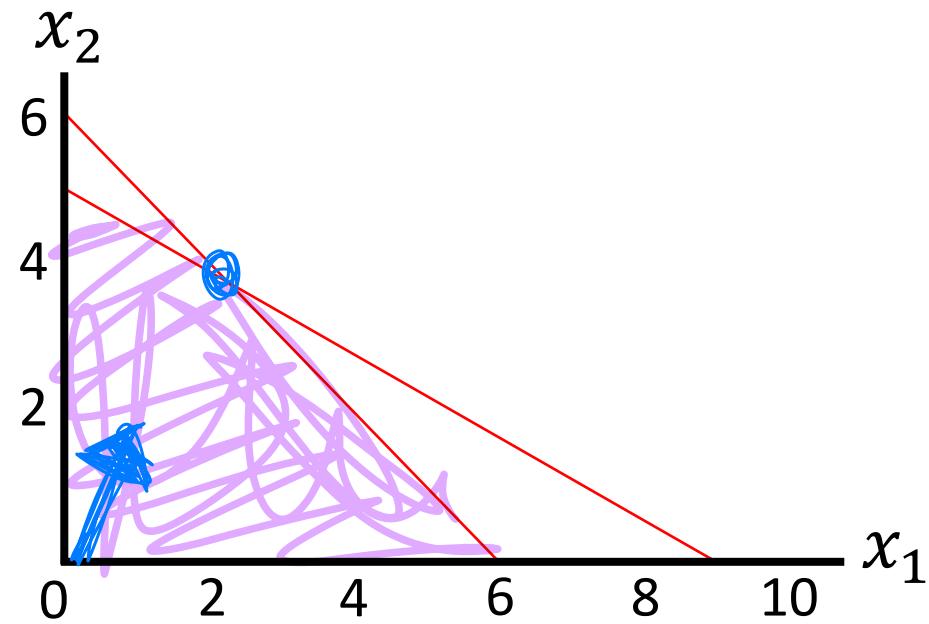
ILP is NP-hard

$x_1, x_2 \in \mathbb{R}$

Objective: $\max 5x_1 + 8x_2$

Subject to: $x_1 + x_2 \leq 6$

 $5x_1 + 9x_2 \leq 45$ $x_1, x_2 \geq 0$

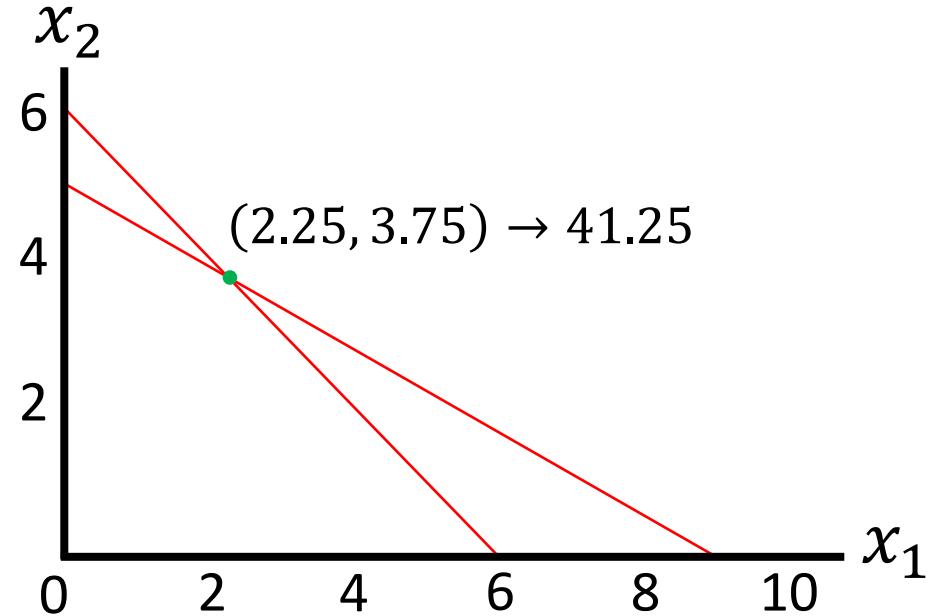


$$x_1, x_2 \in \mathbb{R}$$

Objective: $\max 5x_1 + 8\underline{x}_2$

Subject to:

$$x_1 + \underline{x}_2 \leq 6$$
$$5x_1 + 9x_2 \leq 45$$
$$x_1, x_2 \geq 0$$

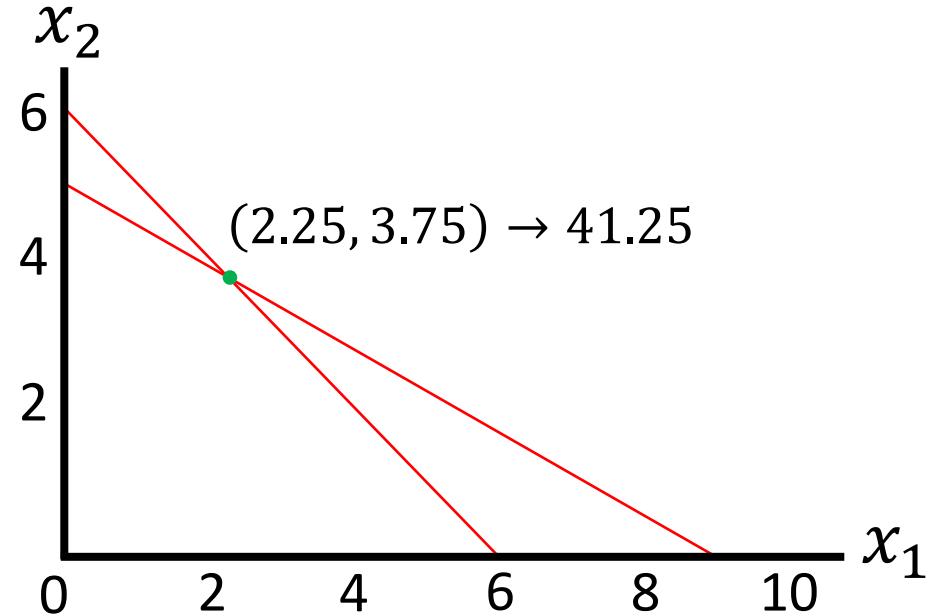


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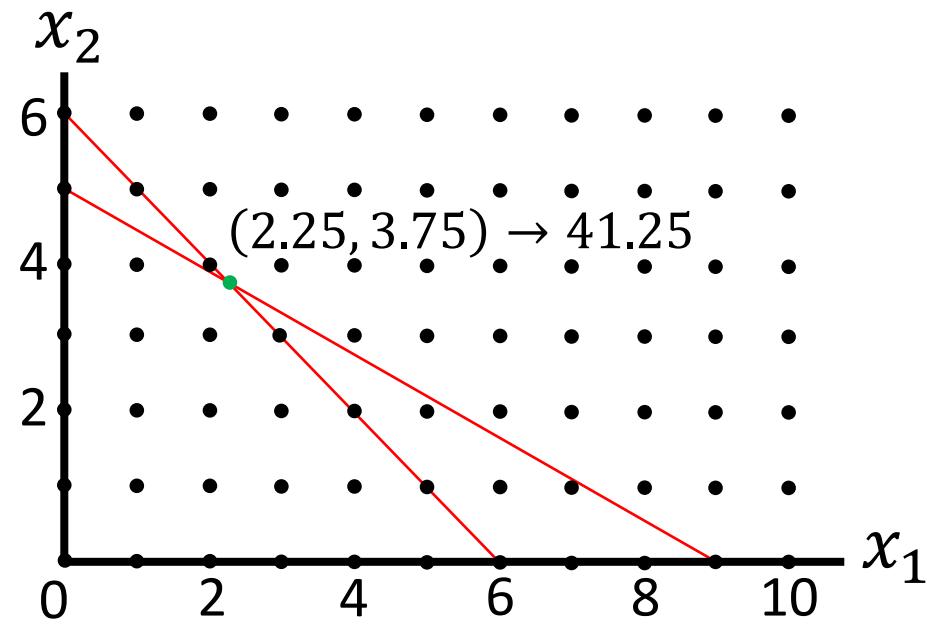


$x_1, x_2 \in \mathbb{N}$ integers ≥ 0

Objective: $\max 5x_1 + 8x_2$

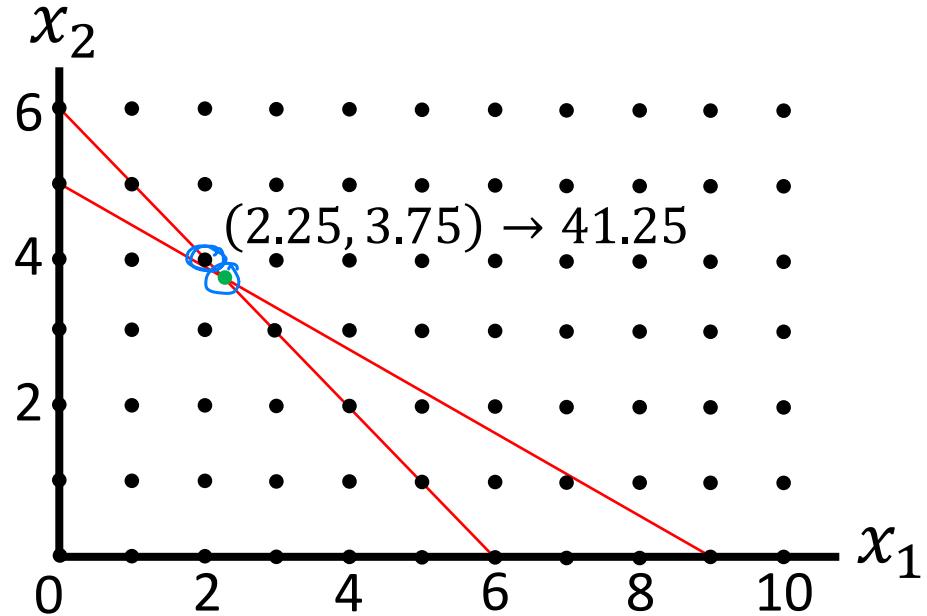
Subject to:

$$x_1 + x_2 \leq 6$$
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$$x_1, x_2 \in \mathbb{N}$$

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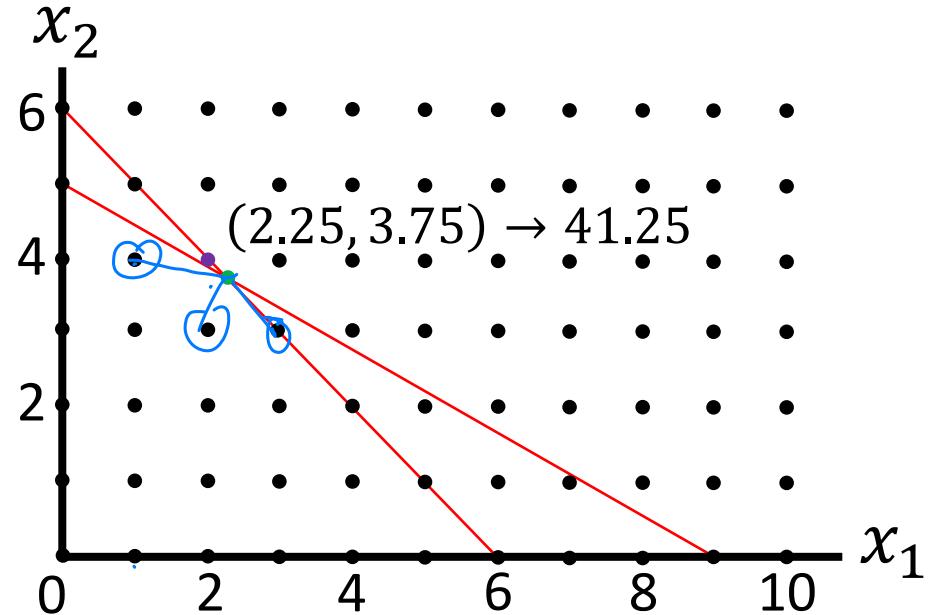
$$x_1, x_2 \in \mathbb{N}$$

$$\text{Objective: } \max 5x_1 + 8x_2$$

$$\begin{aligned} \text{Subject to: } & x_1 + x_2 \leq 6 \\ & 5x_1 + 9x_2 \leq 45 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Optimal continuous solution \rightarrow optimal integer solution?

- Closest integer solution?
- Closest feasible integer solution?
- Closest feasible integer solution on feasible region boundary?



$$x_1, x_2 \in \mathbb{N}$$

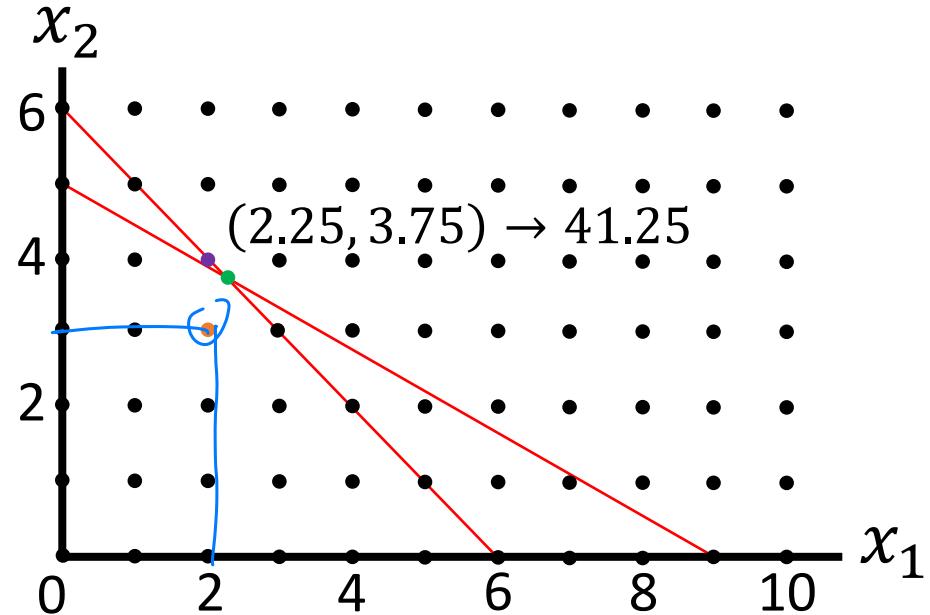
Objective: $\max 5x_1 + 8x_2$

Subject to: $x_1 + x_2 \leq 6$
 $5x_1 + 9x_2 \leq 45$

$$x_1, x_2 \geq 0$$

Optimal continuous solution \rightarrow optimal integer solution?

- Closest integer solution? – Not feasible
- Closest feasible integer solution?
- Closest feasible integer solution on feasible region boundary?



$$x_1, x_2 \in \mathbb{N}$$

Objective:

$$\max 5x_1 + 8x_2$$

Subject to:

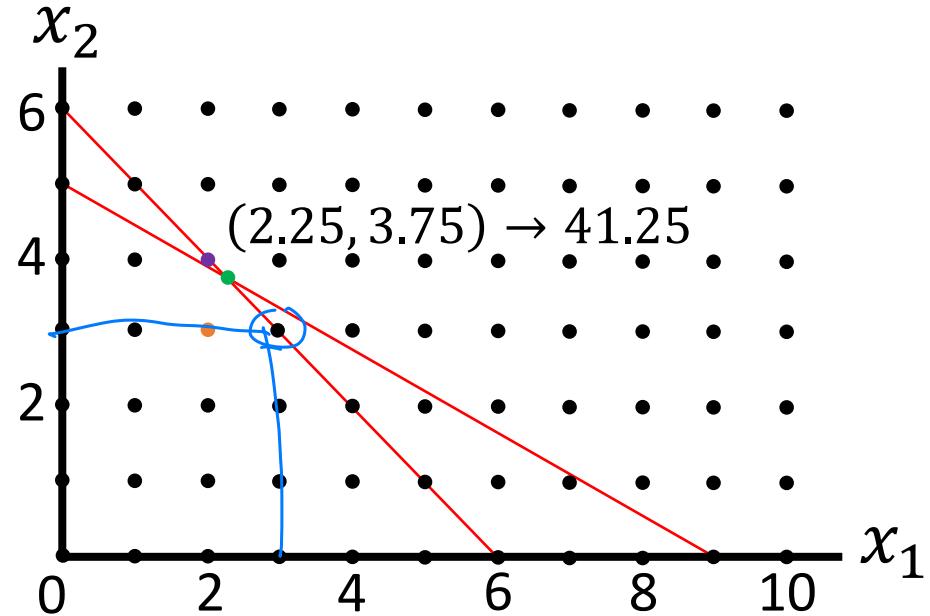
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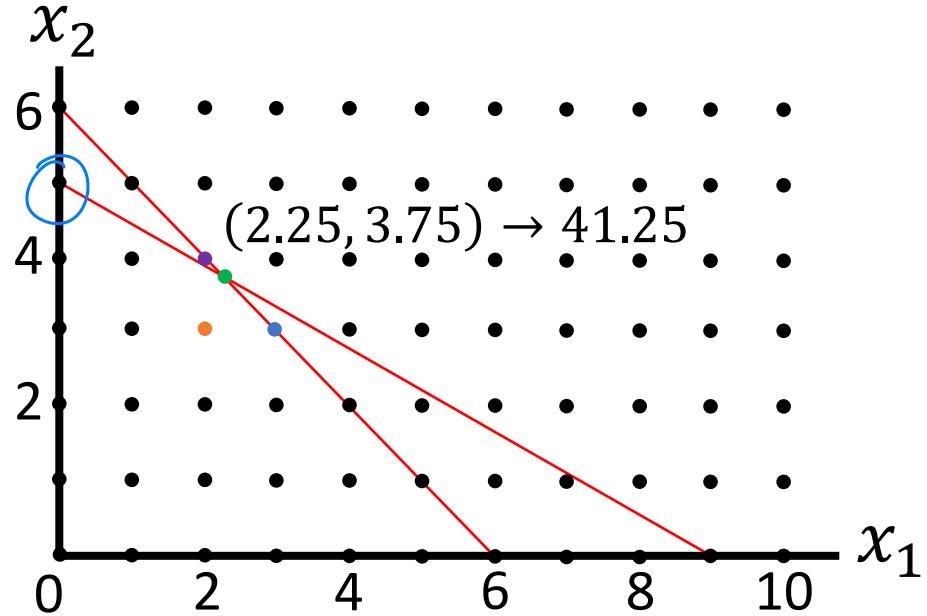
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Optimal continuous solution \rightarrow optimal integer solution?

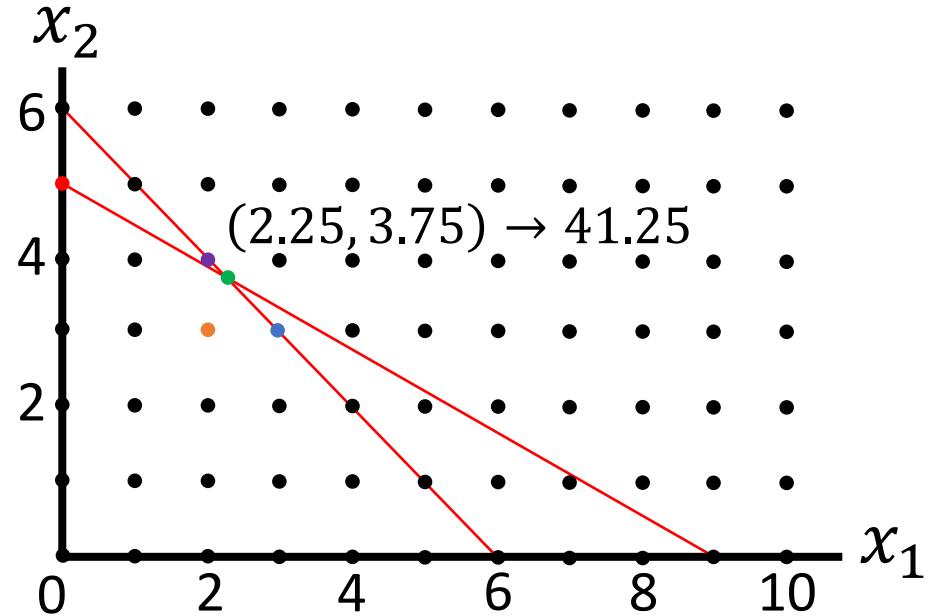
- Closest integer solution? – Not feasible
- Closest feasible integer solution? – Obj = 34
- Closest feasible integer solution on feasible region boundary?



$x_1, x_2 \in \mathbb{N}$	$\textcircled{0}$	\textcircled{S}
Objective:	$\max 5x_1 + 8x_2$	
Subject to:	$x_1 + x_2 \leq 6$	$5x_1 + 9x_2 \leq 45$
	$x_1, x_2 \geq 0$	

Optimal continuous solution \rightarrow optimal integer solution?

- Closest integer solution? – Not feasible
- Closest feasible integer solution? – Obj = 34
- Closest feasible integer solution on feasible region boundary? – Obj = 39



$$x_1, x_2 \in \mathbb{N}$$

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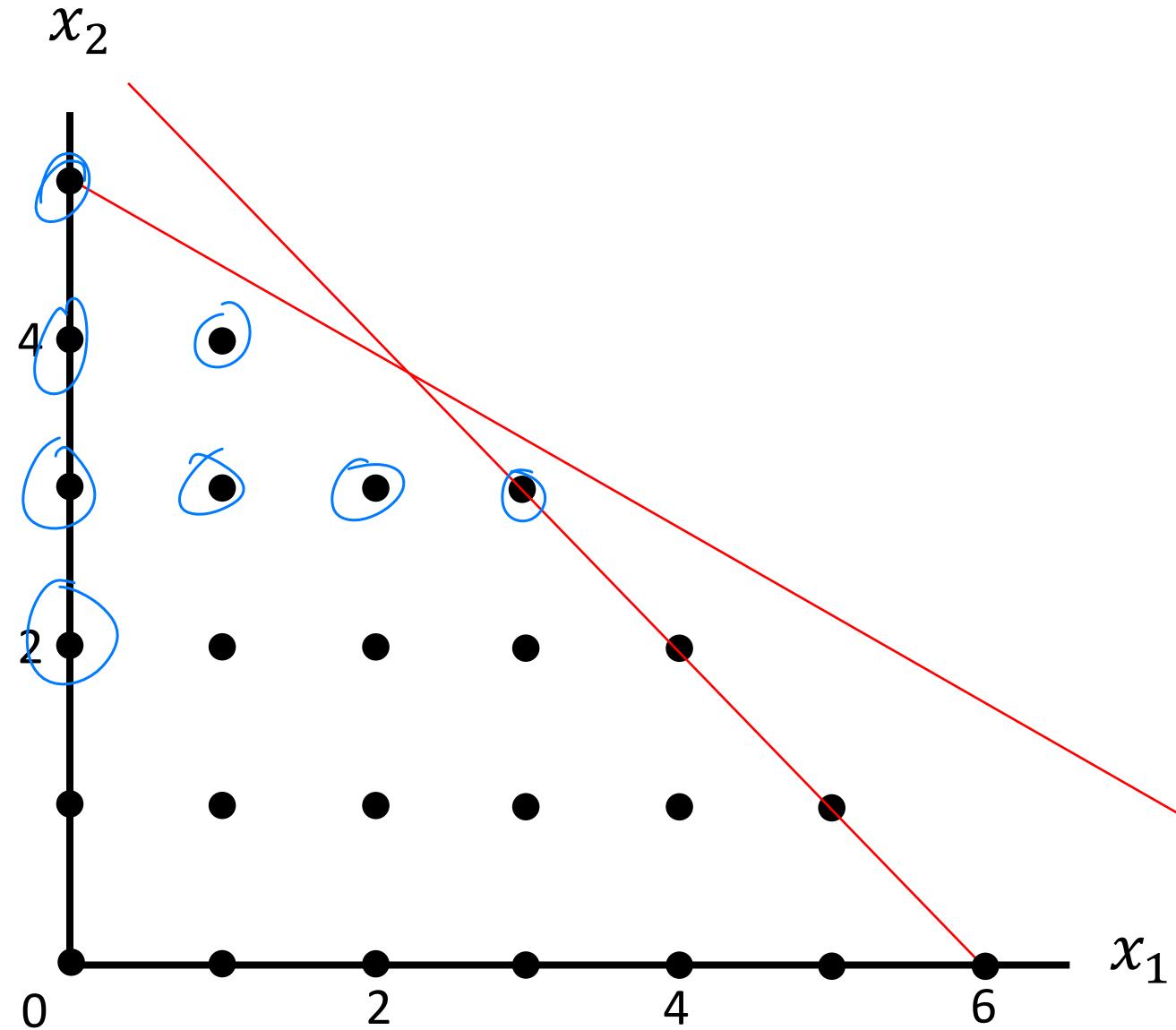
Optimal continuous solution \rightarrow optimal integer solution?

- Closest integer solution? – Not feasible
- Closest feasible integer solution? – Obj = 34
- Closest feasible integer solution on feasible region boundary? – Obj = 39
- **Actual optimal – Obj = 40**

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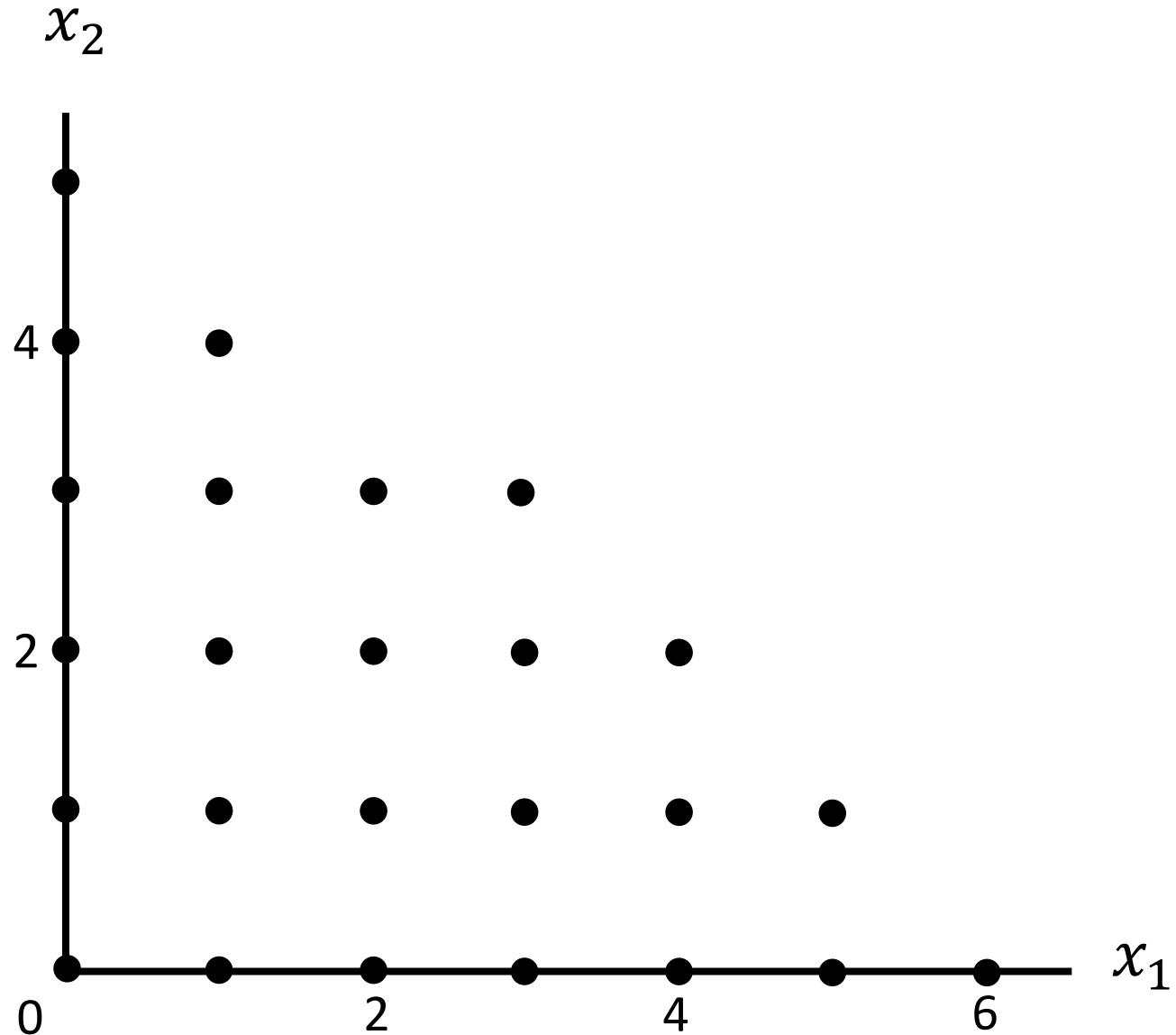
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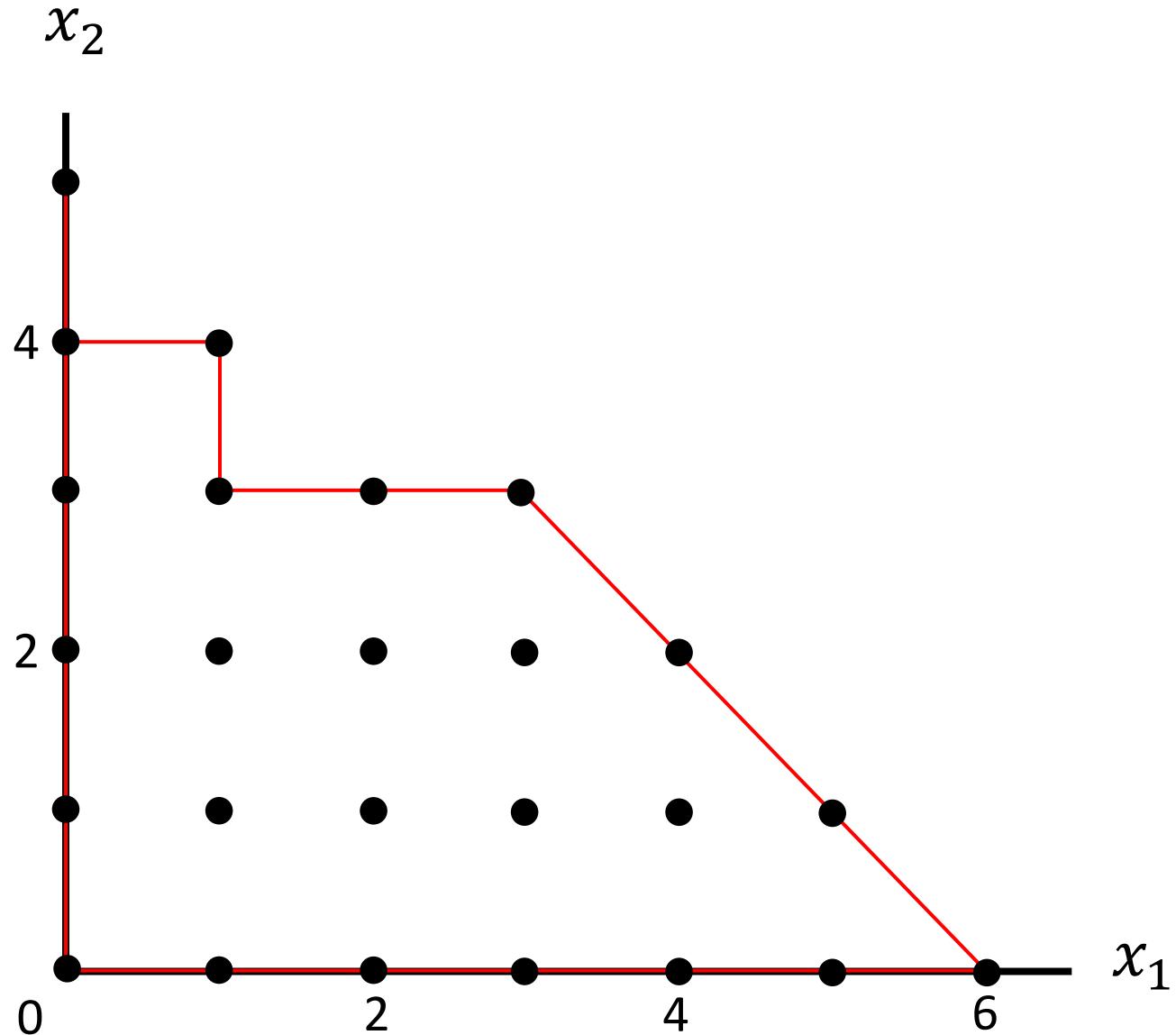


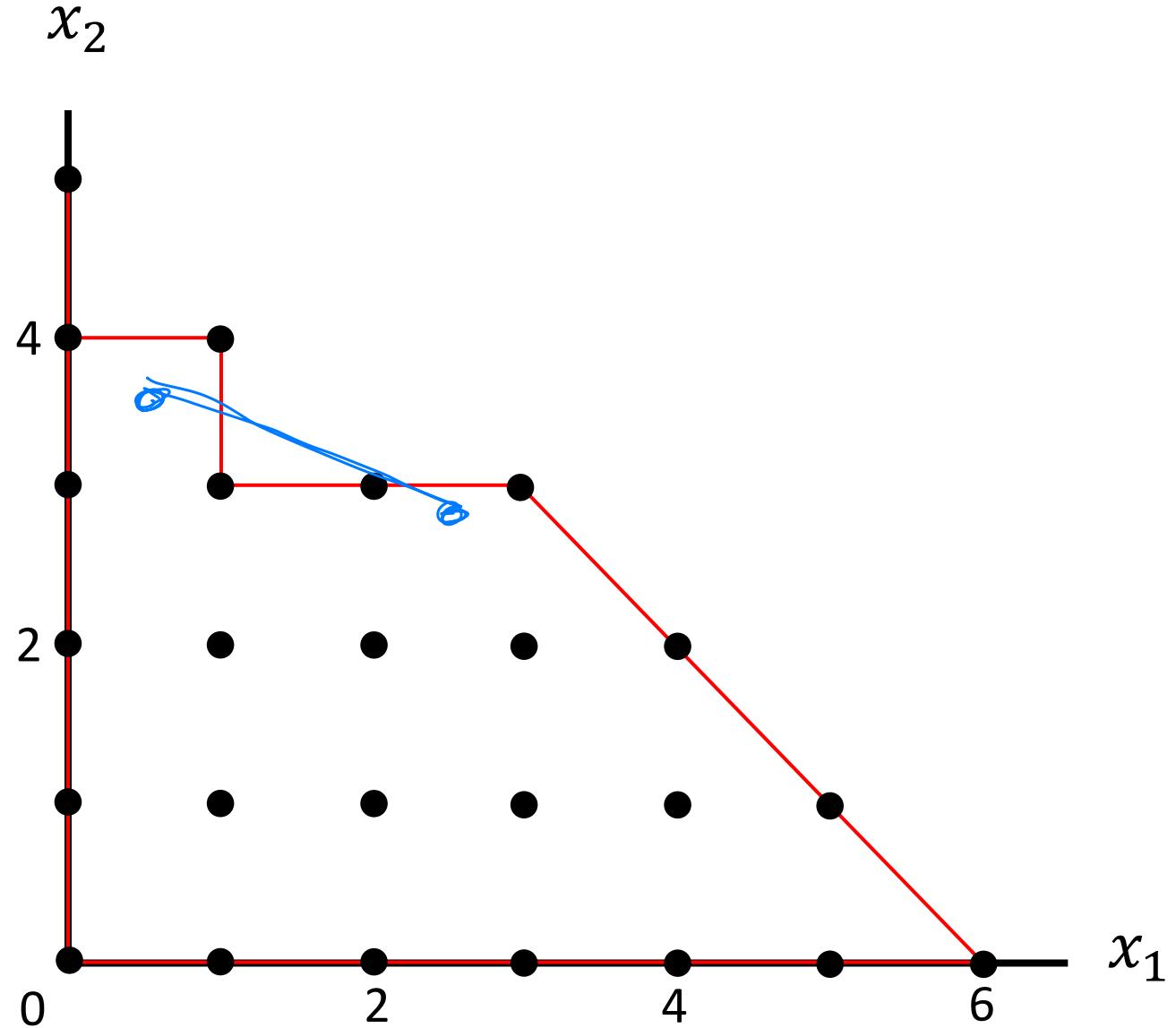
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Integer feasible region:



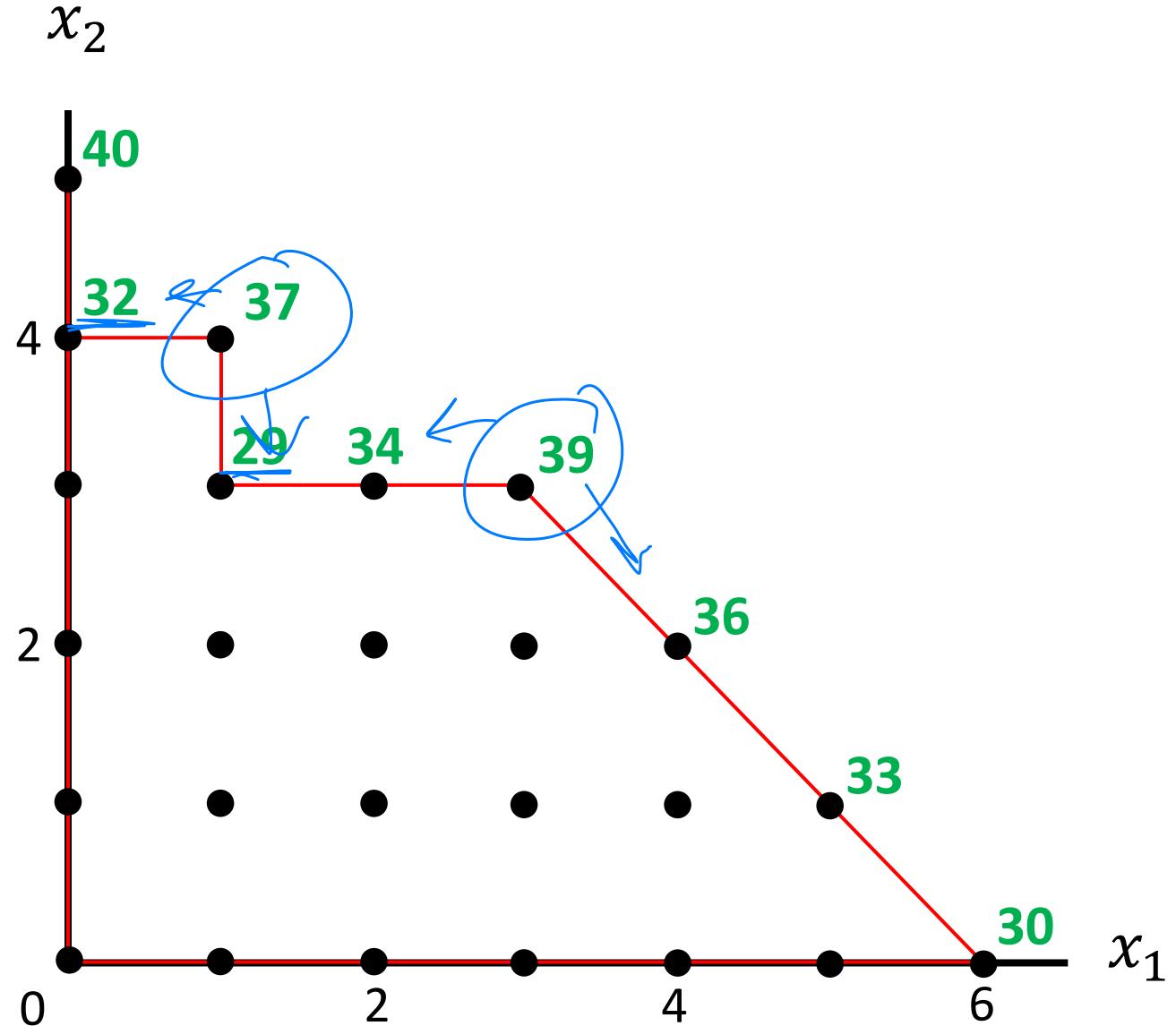


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Integer feasible region:
• Not convex.



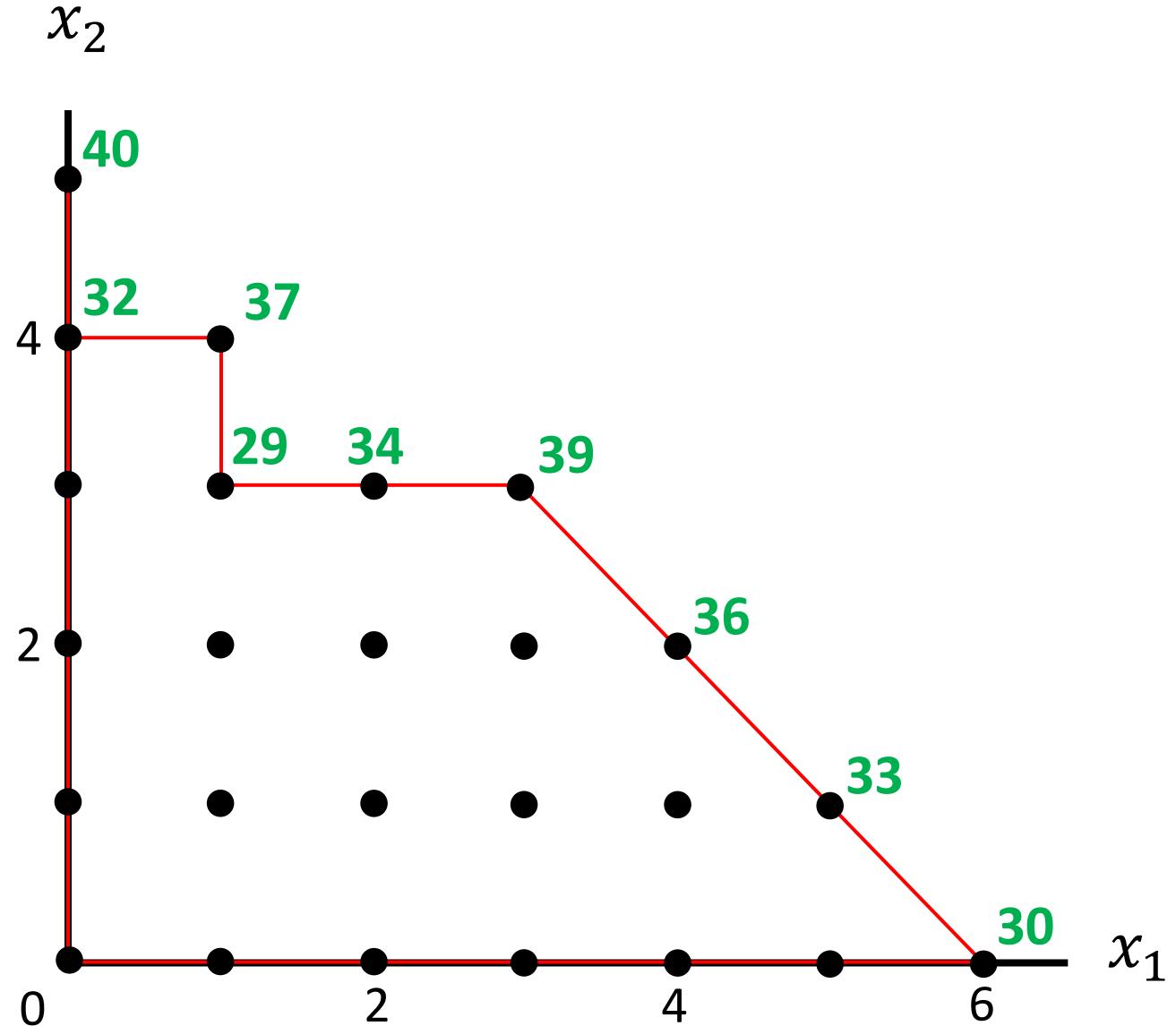
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Integer feasible region:

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$$\begin{aligned} \text{Subject to: } & x_1 + x_2 \leq 6 \\ & 5x_1 + 9x_2 \leq 45 \\ & x_1, x_2 \geq 0 \end{aligned}$$

Integer feasible region:

- Not convex.
- local optimum \neq global optimum.