Given a language of, is mere a TM mat decides it?, is mere If not, f is underidable. Examples of indecidable ranguages: - SELFREJECT = Zem7: M rejects CM73 - SELFHALT = ZEMT: M halts on CM7] - SELFACCEPT = ¿ CM7: Maccept CM7} - ACCEPT = Z<M, w>: Maccepts w] How do we prove a language is undecidable? Prove a language is (1) Diagonalization - cantor proved frat IIIZ/7/Z/ TM Y accept ZY7 <=> TM not accept (2) Reduction from deciding a known undecidable language

Theorem HALT = 2 CM, w7: M halts on w3

is undecidable.

Proof We prove that HALT is underidable by reduction from deciding SELFHALT.

Suppose HALT is decidable. Let H be the TM that decides HALT.

Now we can make a TM called SH mat decides SELFHALT as follows:

TSH(CM7): Lon Hon (M, CM77)



SH is a TM that accepts 2M7 If M halfs on <M) It is a TM that accepts CM, w? iff M halts on W M is a TM that is arbitrary

SH accepts < M7 iff M halts on <a>.

-SH accepts ZM7 if M halts on ZM7:

Mhalts on CM7. Haccepts on CM, CM77. SHaccepts on CM7

-Mhaltson CM7 if SH accepts <M7: contrapositive Strejects curif M does not halt on CM7: p=2qq=7-pM does not halt on CM7. So H rejects CM7. So SH rejects CM7. But SH is underidable! so Il cannot exist.

- MOST languages of form "Griven CM7, does Maccept_

are undeciable

ave there indevidable problems mat are not about This that are indevidable?

1) Diophantine Equations - 1970

Given a polynomial, does it have integer solution?

 $a^{2} + b^{2} = c^{2}$ yes intinitely $a^{n} t b^{n} = c^{n}$ n>2

2) post correspondence Problem